

# Public Key Infrastructure

Yan Huang

Credit: David Evans (UVA)

# Using RSA to Encrypt

- Use 2048-bit modulus (recommended)
- Encrypt 1MB file
  - 4096 2048-bit messages
  - Each  $M^e$  requires  $\log_2(e)$  1024-bit number modular multiplications (unless  $e$  is fixed to some small numbers)
- Why does not one use RSA like this?
  - About 100-1000 times slower than DES
  - Need to be careful not to encrypt particular Ms
  - Can speed up encryption by choosing  $e$  that is an easy number to multiply by (e.g., 3 or  $2^{16} + 1$ )
  - But, decryption must use non-easy  $d$  ( $\sim 2048$  bits)

# Hybrid Encryption

- Use RSA to establish a shared secret key for symmetric cipher (e.g., AES)
  - Encrypt the secret key with OAEP padding
- Sign (encrypt with private key) a hash of the message
  - A short block that is associated with the message

# RSA Paper

“The need for a courier between every pair of users has thus been replaced by the requirement for a single secure meeting between each user and the public file manager when the user joins the system.”

# Key Management

Public keys only useful if you know:

1. The public key matches the entity you think it does (and no one else).
2. The entity is trustworthy.

# Approach 1: Public Announcement

- Publish public keys in a public forum
  - USENET groups
  - Append to email messages
  - New York Time classifieds
- Easy for rogue to pretend to be someone else

# Approach 2: Public Directory

- Trusted authority maintains directory mapping names to public keys
- Entities register public keys with authority in some secure way
- Authority publishes directory
  - Print using watermarked paper, special fonts, etc.
  - Allow secure electronic access

Can we avoid needing an  
on-line directory?

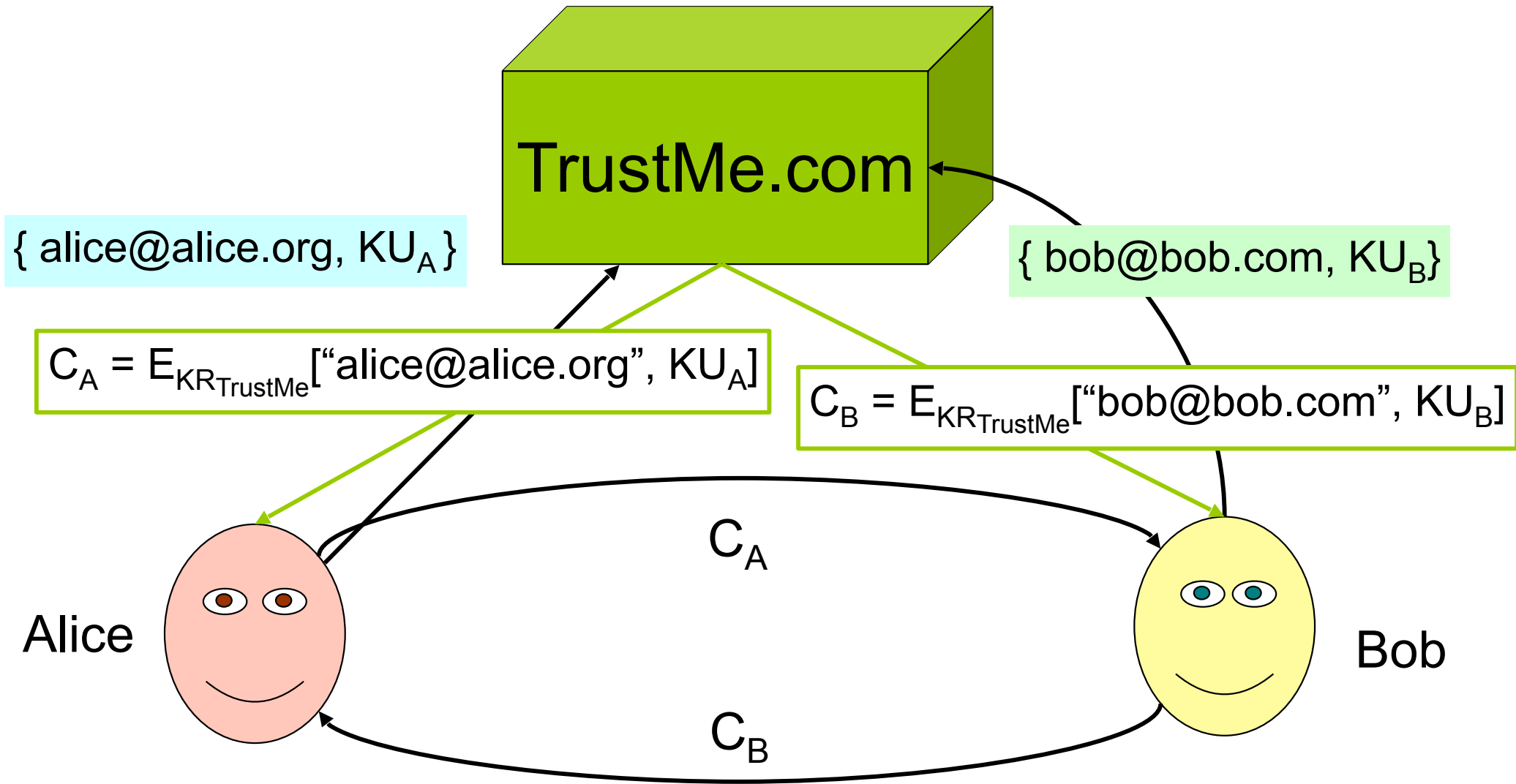


# Certificates

Loren Kohnfelder, MIT 4<sup>th</sup> year thesis project,  
1978: Towards a Practical Public-key  
Cryptosystem

“Public-key communication works best when the encryption functions can reliably be shared among the communicants (by direct contact if possible). Yet when such a reliable exchange of functions is impossible the next best thing is to trust a third party. Diffie and Hellman introduce a central authority known as the Public File... Each individual has a name in the system by which he is referenced in the Public File. Once two communicants have gotten each other’s keys from the Public File then can securely communicate. **The Public File digitally signs all of its transmission so that enemy impersonation of the Public File is precluded.**”

# Certificates



KU — Public Key  
KR — Private Key

### www.google.com

Your connection to this site is private.

Permissions

Connection

#### Cookies

17 from this site, 0 from other sites

#### Permissions

Location: Allowed by you

Microphone: Allowed by you

[Site settings](#)

GeoTrust Global CA  
↳ Google Internet Authority G2  
↳ \*.google.com



#### \*.google.com

Issued by: Google Internet Authority G2

Expires: Tuesday, May 3, 2016 at 8:00:00 PM Eastern Daylight Time

✓ This certificate is valid

#### Details

Subject Name	
Country	US
State/Province	California
Locality	Mountain View
Organization	Google Inc
Common Name	*.google.com
Issuer Name	
Country	US
Organization	Google Inc
Common Name	Google Internet Authority G2
Serial Number	1403480822151576483
Version	3
Signature Algorithm	SHA-256 with RSA Encryption ( 1.2.840.113549.1.1.11 )
Parameters	none
Not Valid Before	Thursday, February 4, 2016 at 5:22:52 AM Eastern Standard Time
Not Valid After	Tuesday, May 3, 2016 at 8:00:00 PM Eastern Daylight Time

OK

Data encrypted using secret key exchanged using some public key associated with some certificate.

GeoTrust Global CA  
↳ Google Internet Authority G2  
↳ \*.google.com

Parameters none

**Not Valid Before** Thursday, February 4, 2016 at 5:22:52 AM Eastern Standard Time

**Not Valid After** Tuesday, May 3, 2016 at 8:00:00 PM Eastern Daylight Time

**Public Key Info**

**Algorithm** Elliptic Curve Public Key ( 1.2.840.10045.2.1 )

**Parameters** Elliptic Curve secp256r1 ( 1.2.840.10045.3.1.7 )

**Public Key** 65 bytes : 04 D4 1A 7D E1 68 CB BC 89 4A 30 82 49  
0C 90 38 BC 1A FC 8E EB 68 42 45 6D 08 E9 2B 84  
27 B6 D0 67 08 A4 21 63 FA D5 21 0A F8 8C 49 76  
CE 04 6A F9 F9 02 07 8A D8 E5 24 8B 3F AC 93 32  
D8 97 DE EA

**Key Size** 256 bits

**Key Usage** Encrypt, Verify, Derive

**Signature** 256 bytes : 1E 14 BB FC 57 B8 E6 16 ...

**Extension** Key Usage ( 2.5.29.15 )

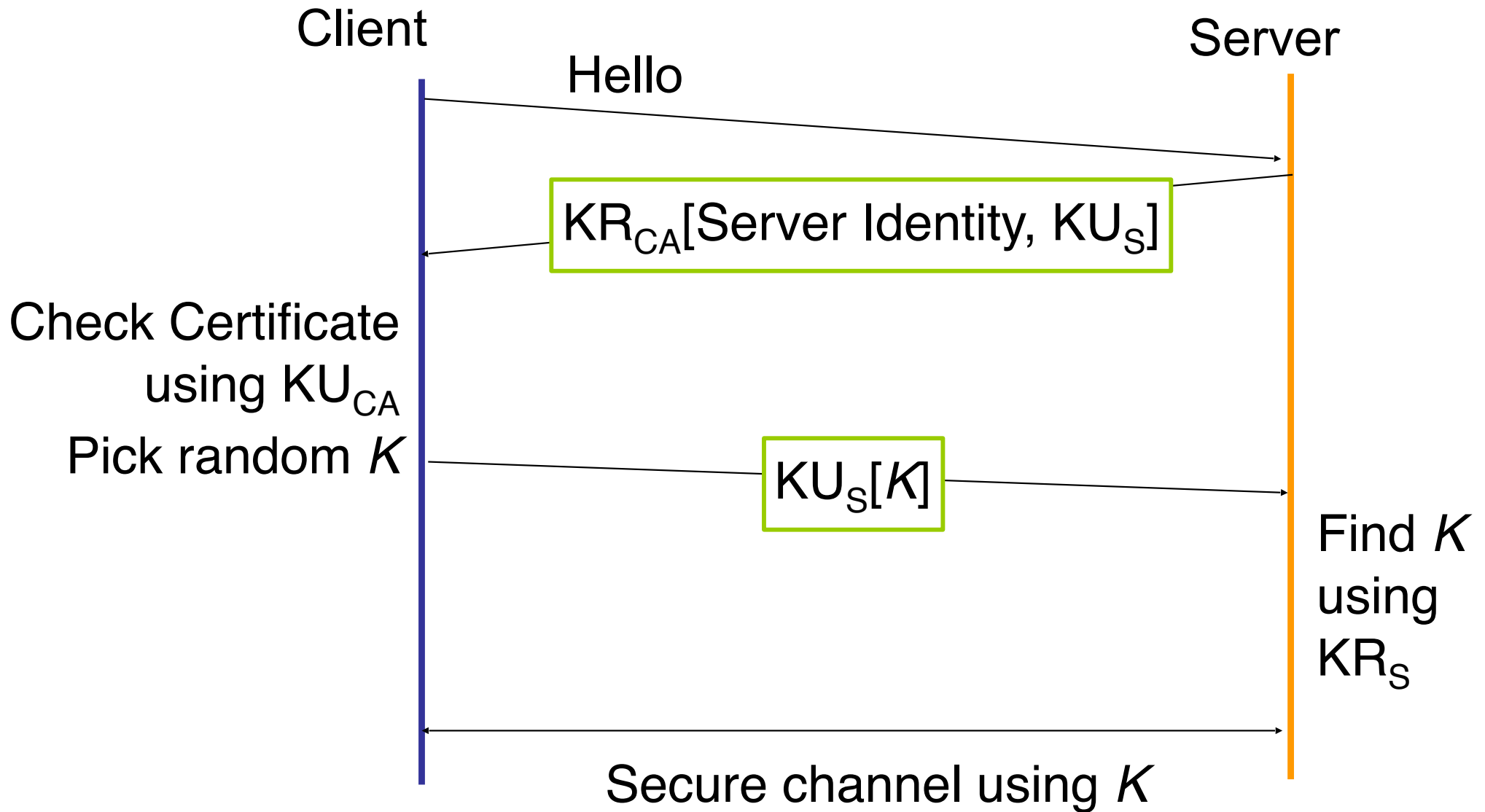
**Critical** NO

**Usage** Digital Signature

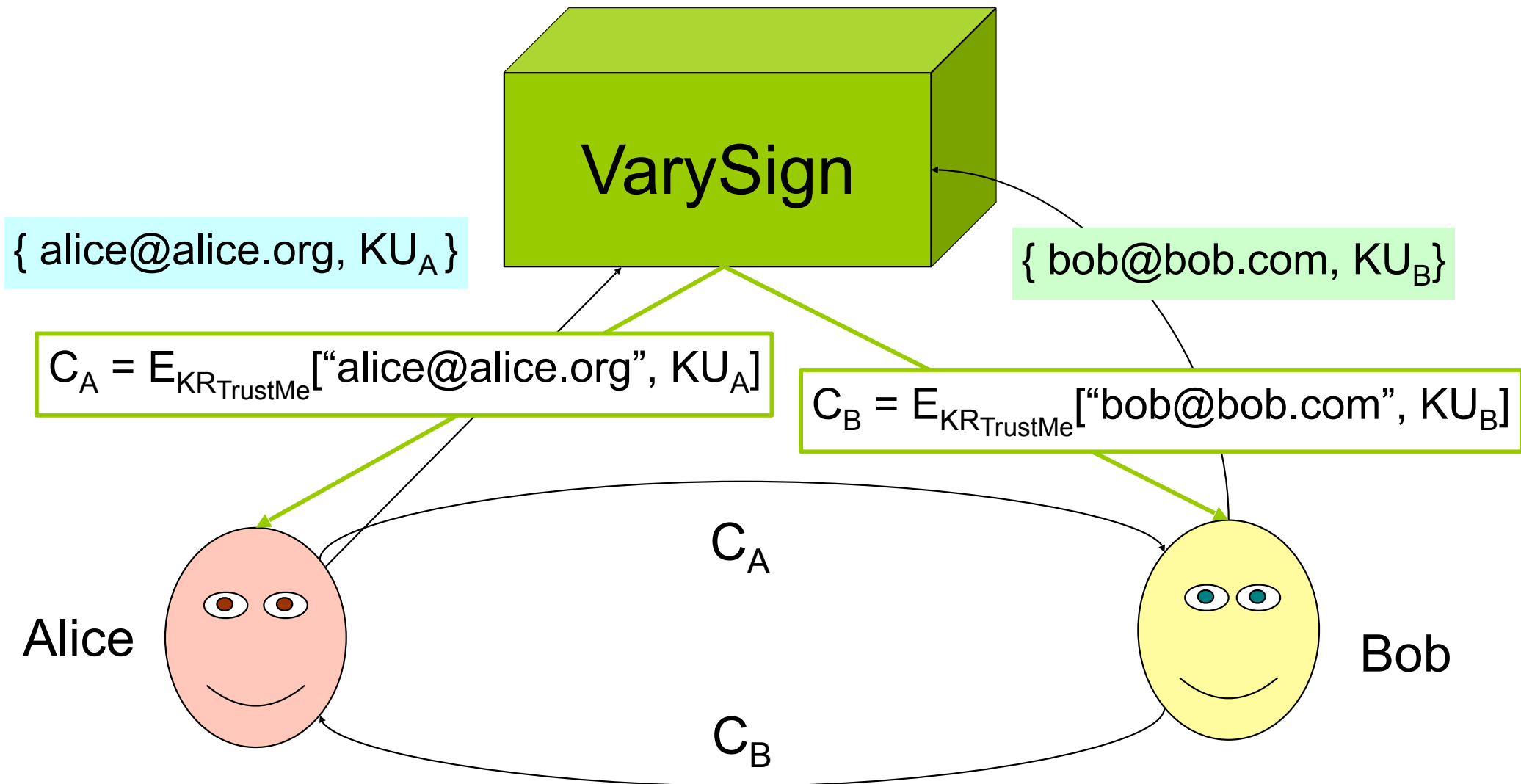
OK

# SSL (Secure Sockets Layer)

## Simplified TLS Handshake Protocol

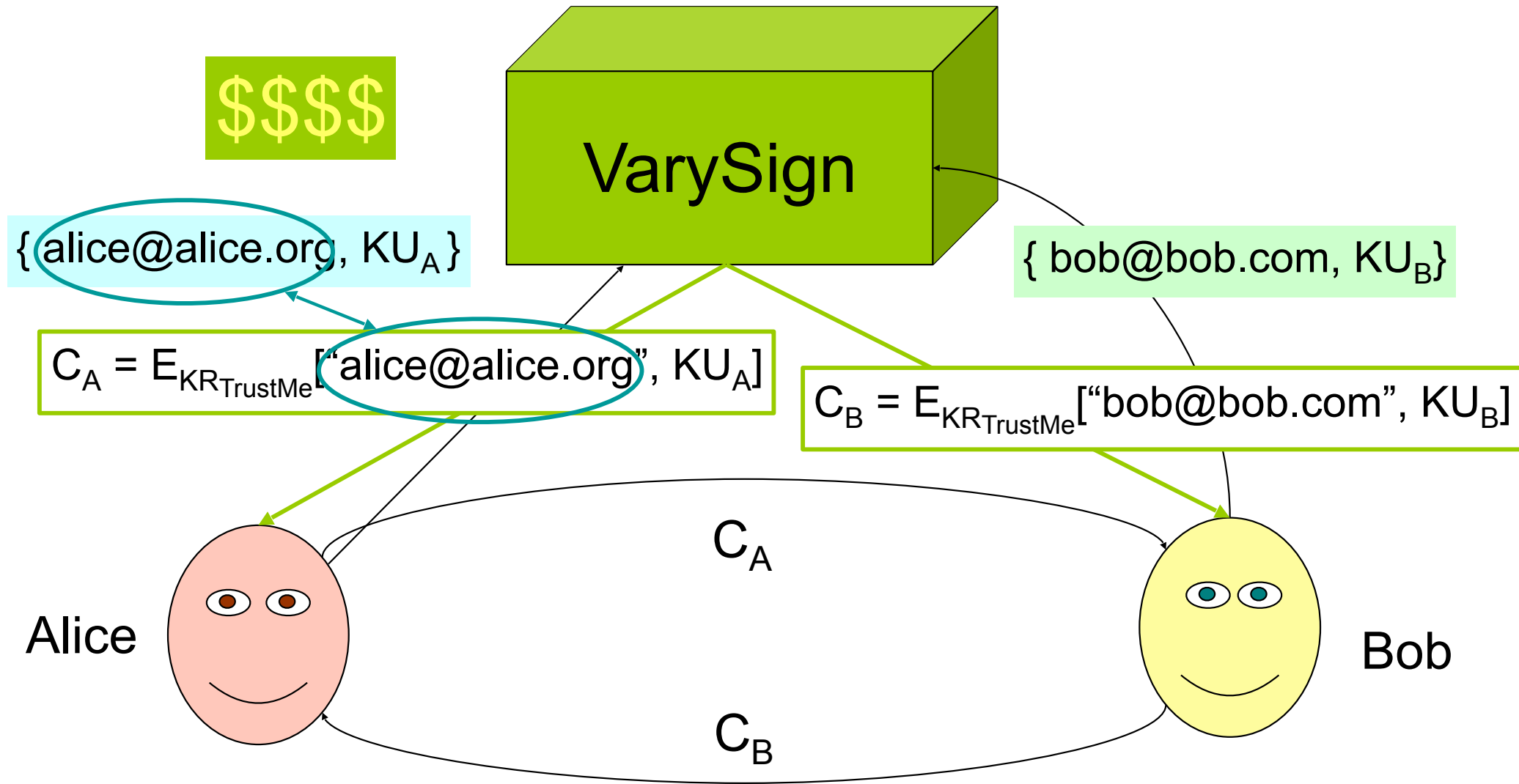


# Certificates



How does TrustMe.com decide whether to provide Certificate?

# Verifying Identities





Compare All SSL Certificates

Close Window

Option	Secure Site SSL Certificates	Secure Site Pro True 128-Bit SSL	Commerce Site SSL Certificates	Commerce Site Pro True 128-bit SSL	Managed PKI for SSL Standard Edition	Managed PKI for SSL Premium Edition	Managed PKI for Intranet SSL Standard Edition	Managed PKI for Intranet SSL Premium Edition
View Product Description	<a href="#">Product Info</a>	<a href="#">Product Info</a>	<a href="#">Product Info</a>	<a href="#">Product Info</a>	<a href="#">Product Info</a>	<a href="#">Product Info</a>	<a href="#">Product Info</a>	<a href="#">Product Info</a>
Ready to Buy?	<a href="#">Buy Now</a>	<a href="#">Buy Now</a>	<a href="#">Buy Now</a>	<a href="#">Buy Now</a>	<a href="#">Contact Sales</a>	<a href="#">Contact Sales</a>	<a href="#">Contact Sales</a>	<a href="#">Contact Sales</a>
Price: 2-Year Certificate	\$598	\$1,790	\$1,798	\$2,795	Contact Sales	Contact Sales	Contact Sales	Contact Sales
Price: 1-Year Certificate	\$349	\$995	\$949	\$1,495	\$249/certificate for 10	\$695/certificate for 10	\$179/certificate for 10	\$489/certificate for 10
Number of certificates	Single	Single	Single	Single	5, 10, 25, 50, 100, more	5, 10, 25, 50, 100, more	5, 10, 25, 50, 100, more	5, 10, 25, 50, 100, more
Free SSL Trial	Free SSL Trial	-	-	-	-	-	-	-
Minimum SSL Encryption	40-bit	128-bit	-	-	-	-	-	-
Issuance	Standard	Express delivery	-	-	-	-	-	-
Online Payment Processing	-	-	-	-	-	-	-	-
VeriSign NetSure Protection Warranty	\$100,000	\$250,000	-	-	-	-	-	-
Qualys Guard	-	Free	-	-	-	-	-	-
Netcraft	-	Free \$1,000 value	Free \$1,000 value	Free \$1,000 value	-	-	-	-
Keynote Red Alert	-	Free	Free	Free	-	-	-	-
VeriSign Secured Seal								
Authentication	2 factor authentication	2 factor authentication	2 factor authentication	2 factor authentication	Class 3 organizational authentication	Class 3 organizational authentication	Class 3 organizational authentication	Class 3 organizational authentication

With over half a million businesses authenticated, VeriSign follows a rigorous and independently audited authentication process. All involved VeriSign employees pass stringent background checks, and each authentication is split between multiple individuals. We maintain physically secure facilities, including biometric screening on entry.



# VeriSign's Certificate Classes

- “Secure Site” SSL Certificate
  - Supports 40-bit session key
  - Proves: you are communicating with someone willing to pay VeriSign \$598 (or with ~\$1000 to break a 40-bit key)
  - Except they have a free 14-day trial (but it uses a different Trial CA key)



## Why SGC-Enabled Certificates?

Without an SGC-enabled certificate in place, site visitors using certain older browsers and many Windows 2000 users will only receive 40- or 56-bit encryption.

VeriSign (including its affiliates, resellers, and subsidiaries) is the only leading SSL provider with SGC-enabled SSL Certificates that can provide true 128- or 256-bit encryption to the most Web site visitors.

128-bit encryption offers  $2^{88}$  times as many possible combinations as 40-bit encryption. That's over a trillion times a trillion times stronger.

**Isn't your customer's trust worth it?**

Close Window

# “Secure Site Pro” Certificate

- \$995 per year
- “true 128-bit key”
  - “128-bit encryption offers  $2^{88}$  times as many possible combinations as 40-bit encryption. That’s over a trillion times a trillion times stronger.”
    - $\text{trillion} = 10^{12}$        $\text{trillion} * \text{trillion} = 10^{24}$
  - Verisign’s marketing claim could be:
    - “trillion times a trillion times a trillion times a trillion times a trillion times a trillion times a trillion times ten thousand (in Britain it is a trillions time a trillion times a trillion times a trillion times a billion times a thousand) times stronger”
      - (but that would sound even sillier!)
- Businesses authentication: “out-of-band” communication, records

VeriSign Inc. - www.verisign.com - Microsoft Internet Explorer

File Edit View Favorites Tools Help

Back Forward Stop Home Search Favorites Media Print Mail News Chat

Address <http://www.verisign.com/developer/notice/authenticode/> Go Links >>



HOME | PRODUCTS | CONSULTING | TRAINING | SUPPORT | CORPORATE

**SEARCH:**

All VeriSign

**CORPORATE**

- Hot Jobs
- Investor Relations
- News Center
- Contact Us

**RESOURCES**

- Repository
- Site Map

**VERISIGN**

International

[Home](#) > Security Notice

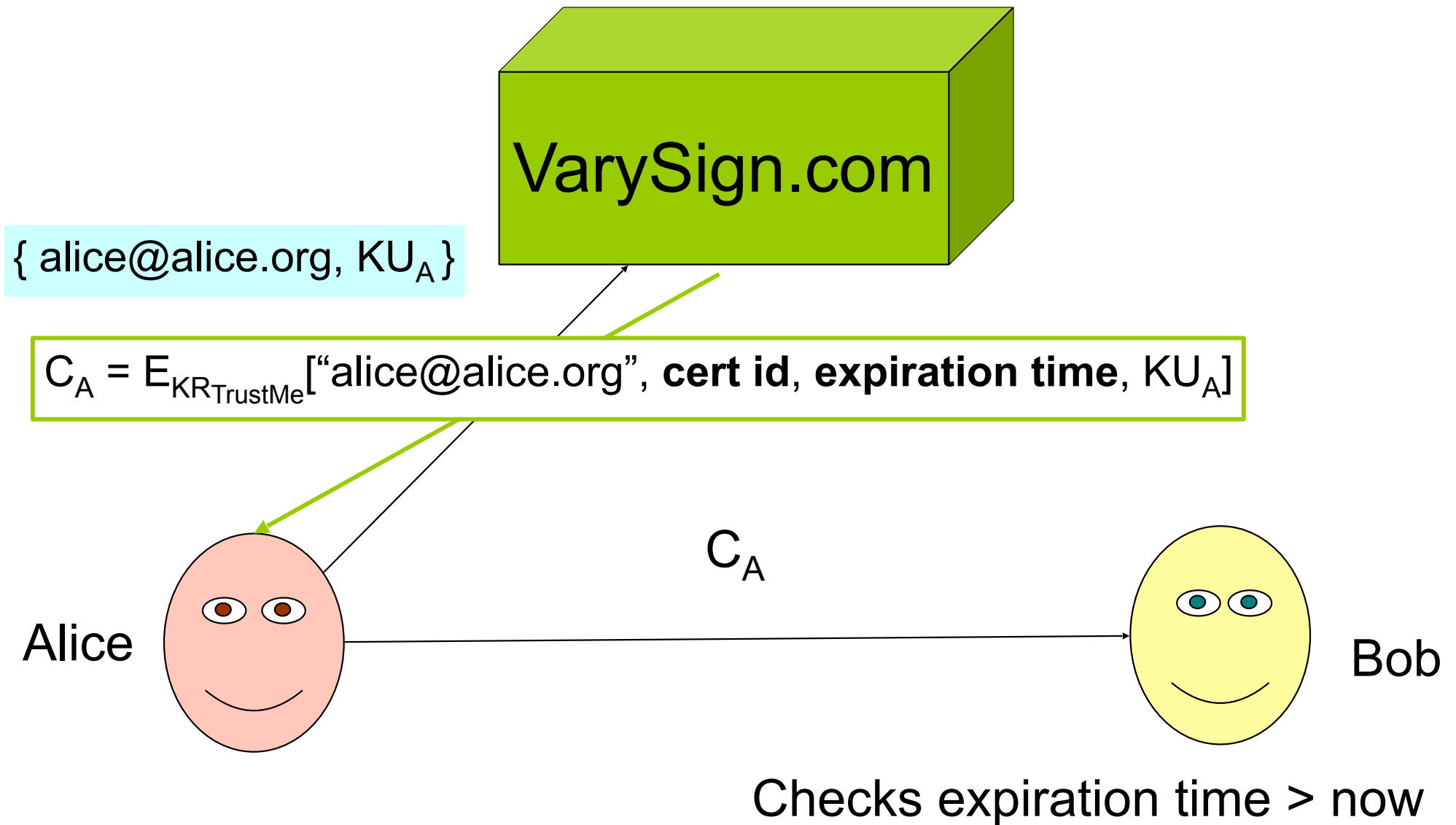
## VeriSign Security Alert Fraud Detected in Authenticode Code Signing Certificates March 22, 2001

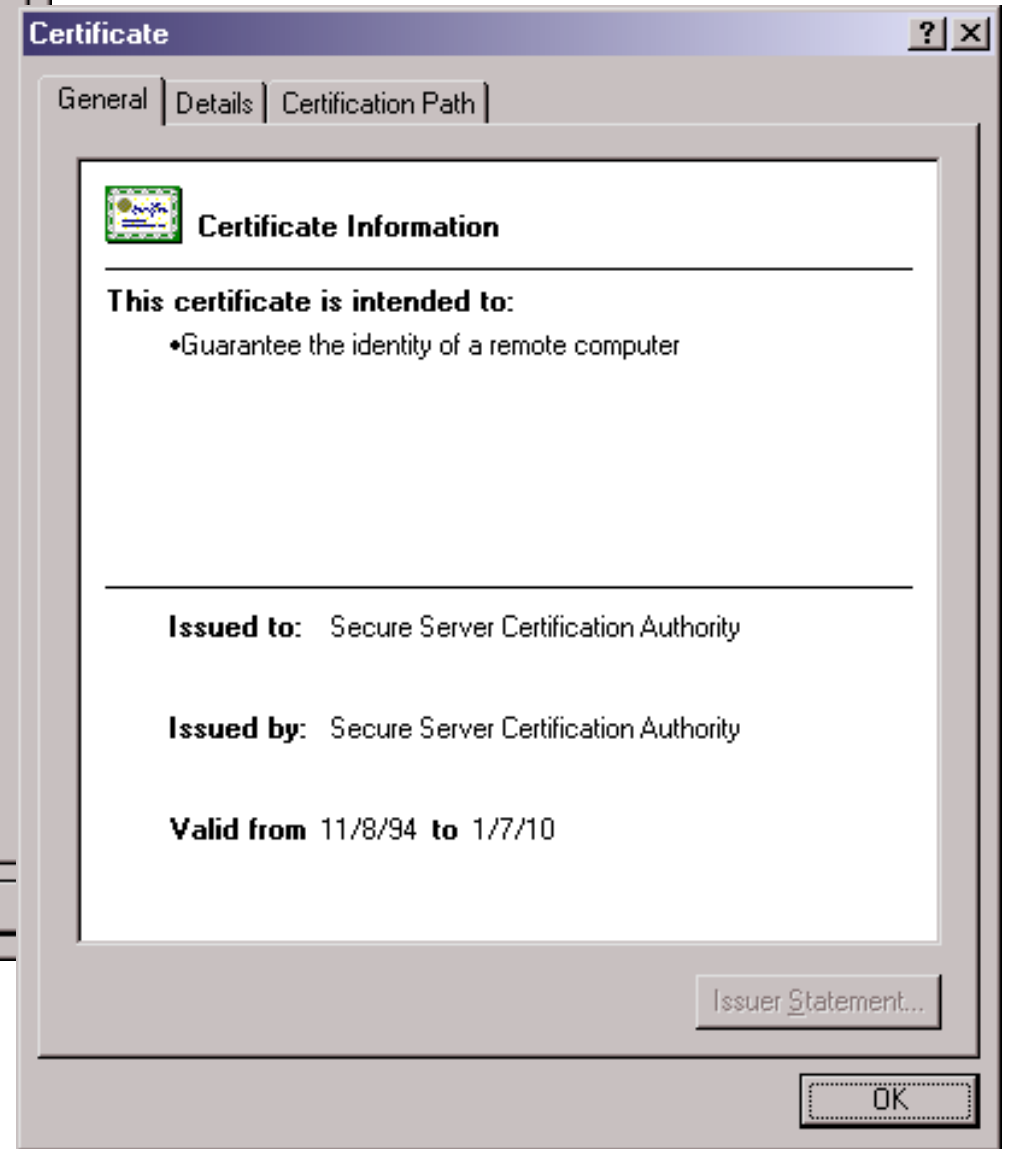
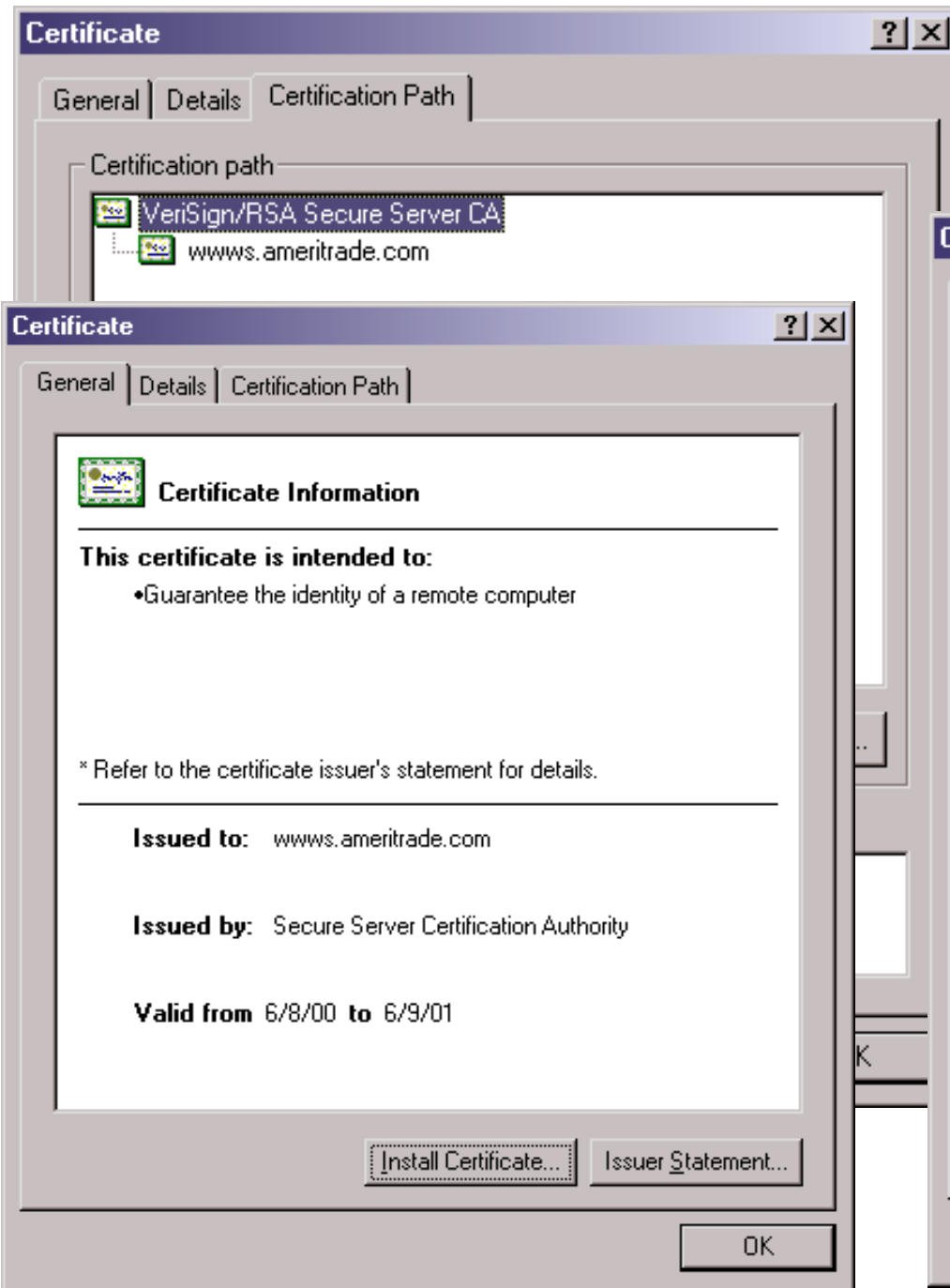
VeriSign, Inc, discovered through its routine fraud screening procedures that on 29 and 30 January 2001, it issued two digital certificates to an individual who fraudulently claimed to be a representative of Microsoft Corporation. VeriSign immediately revoked the certificates. The updated certificate revocation list (CRL) is available at <http://crl.verisign.com/Class3SoftwarePublishers.crl> or through VeriSign real-time Online Certificate Status Protocol (OCSP) Services.

The certificates were VeriSign Class 3 Software Publisher certificates and could be used to sign executable content under the name "Microsoft Corporation". The risk associated with these certificates is that the fraudulent party could produce digitally signed code and appear to be Microsoft Corporation. In this scenario, it is possible that the fraudulent party could create a destructive program or ActiveX control, then sign it using either certificate and host it on a Web site or distribute it to other Web sites.

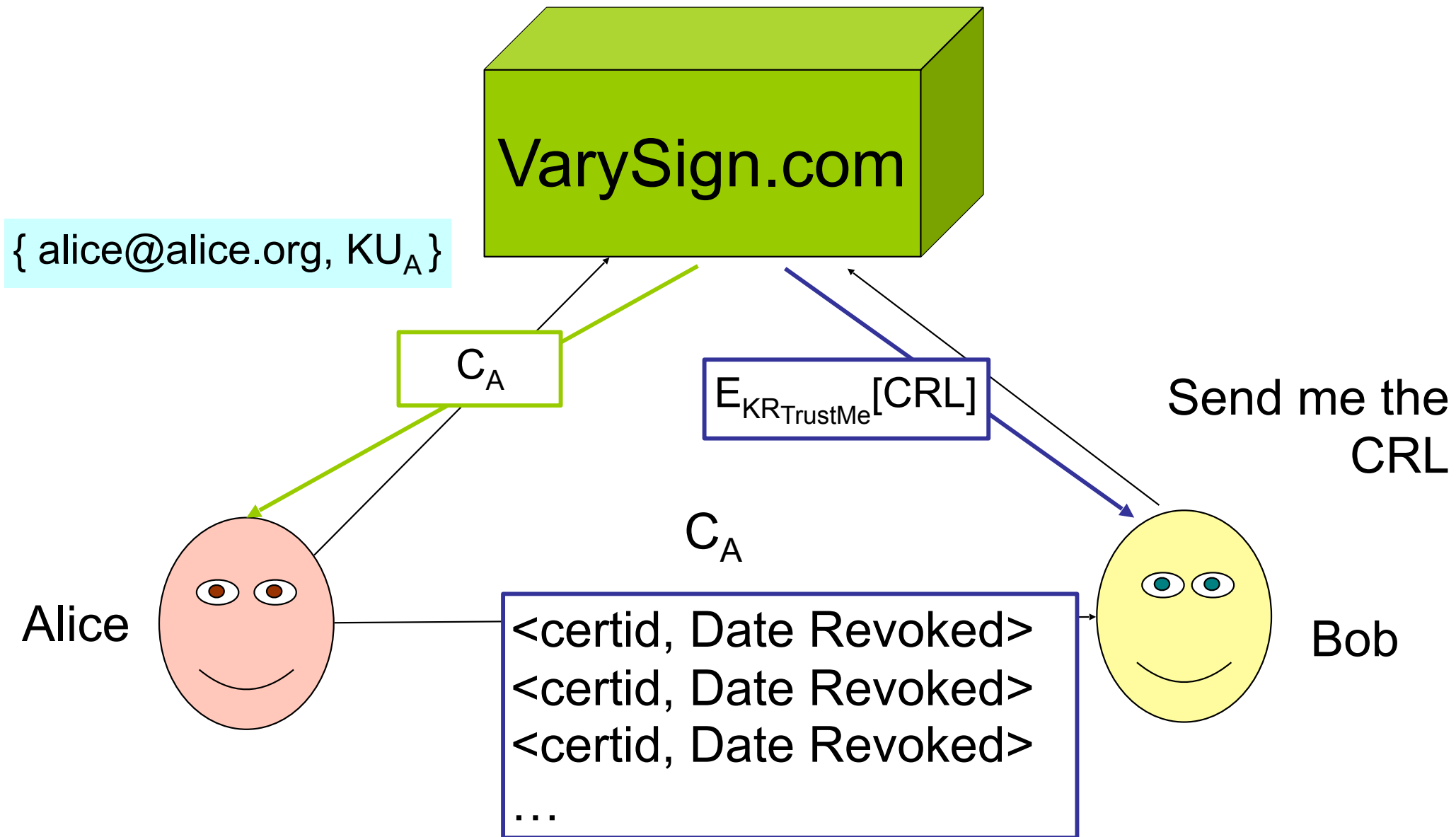
Local intranet

# Limiting The Damage





# Revoking Certificates



# Certificate Questions

- How do participants acquire the authority's public key?
- If authority's private key is compromised, everything is vulnerable!
  - Keep the key locked up well



# Problems with Certificates

- Depends on a certificate authority
  - Needs to be a big, trusted entity
  - Needs to make money (or be publicly funded)
- Need to acquire a certificate
  - Makes anonymity difficult
  - Requires handshaking

# PGP (Pretty Good Privacy)

- Keyring: list of public keys, signed by owner's private key

Alice's keyring:

$E_{KR_{Alice}} (< \text{"bob@bob.com"}, KU_{Bob} >, < \text{"cathy@sharky.com"}, KU_{Cathy} >)$

- Exchanging Keyrings (Web of Trust)
  - Complete Trust: I trust Alice's keyring (add the public key pairings to my own keyring)
  - Partial Trust: I sort of trust Alice, but require confirmation from someone else too (I need to get  $E_{KRCathy} (< \text{"bob@bob.com"}, KU_{Bob} >)$  before trusting  $KU_{Bob}$ )

# Charge

- Look at the certificate chains when you browse the web
  - Find a certificate with a trust chain more than two levels deep
- Update your browser CRLs: when were they last updated?

# Avoiding Certificates

- What if your identity (e.g., your email address) is your public key?
- Is it possible to do this with RSA?

Do you want your email address to be a 200-digit “random” number?

# Identity Based Encryption

- [Shamir 1984], [Boneh & Franklin 2003]

public-key = identity

private-key =  $F(\text{master-key}, \text{identity})$

The owner of the master-key is the new authority. Must be careful who it gives private keys to.

# Key-Generating Service

- Holds master-key
- Participants request private keys from KGS
  - Sends  $s$  to KGS, requests corresponding private key
  - KGS authenticates requestor
  - If valid, computes  $F(\text{master-key}, s)$  and sends over secure channel
- How does the trust given to the KGS compare to that given to CA in SSL?

KGS can decrypt all messages! With certificates, certificate owner still has her own private key.

But, CA can impersonate anyone by generating a certificate with a chosen public-key.

# Shamir's IBE Signature Scheme

- Setup: done by KGS
  - Select  $p, q$  large primes
  - $N = pq$
  - Choose  $e$  relatively prime to  $\varphi(N) = (p-1)(q-1)$
  - Choose  $d$  satisfying  $ed \equiv 1 \pmod{\varphi(N)}$
  - Choose  $h$  a cryptographic hash function
- Publish  $N, e$  and  $h$  to all participants
- Keep  $d$  secret master-key

# Shamir's Signatures

- Generating a private key
$$\text{privatekey}(ID) = ID^d \bmod N$$
  - Can only be done by KGS ( $d$  is master secret)
- Signing a message  $M$  with identity  $ID$ 
  - Obtain  $g = \text{privatekey}(ID)$  from KGS
  - Choose random  $r$  less than  $N$
  - Compute signature  $(s, t)$ :
$$t = r^e \bmod N$$
$$s = g r^{h(t \parallel M)} \bmod N$$

Warning: book typesetting is off and wrong range for  $h$ !



# Verifying a Signature

- KGS produced  $g = ID^d \bmod N$
- Recipient knows  $ID$  and  $M$ , system parameters  $e$  and  $N$

$$t = r^e \bmod N$$

$$s = g r^{h(t \parallel M)} \bmod N \quad (ID^d r^{h(t \parallel M)})^e \bmod N$$

- Verify  $(ID, s, t, M)$

$$\equiv ID^{de} r^{h(t \parallel M)e} \bmod N$$

$$\equiv ID r^{eh(t \parallel M)} \bmod N$$

$$s^e = ID t^{h(t \parallel M)} \bmod N \quad \equiv ID t^{h(t \parallel M)} \bmod N$$

What does non-forgibility of a Shamir IBE signature rely on?

# Identity-Based Encryption

- Shamir's scheme – signatures only, not encryption
- Boneh & Franklin, 2001
  - First practical, provably secure IBE scheme
  - Builds on elliptic curves

# Issues in IBE

- Complete trust in KGS
  - With Boneh & Franklin's system can use secret sharing techniques to divide this trust among multiple entities
  - Could you do this with Shamir's IBE signatures?
- Revocation
  - Can include expiration times in identities
  - But no way to revoke granted private keys