

# Authenticated Encryption

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Credit: Dan Boneh (Stanford, Crypto I)

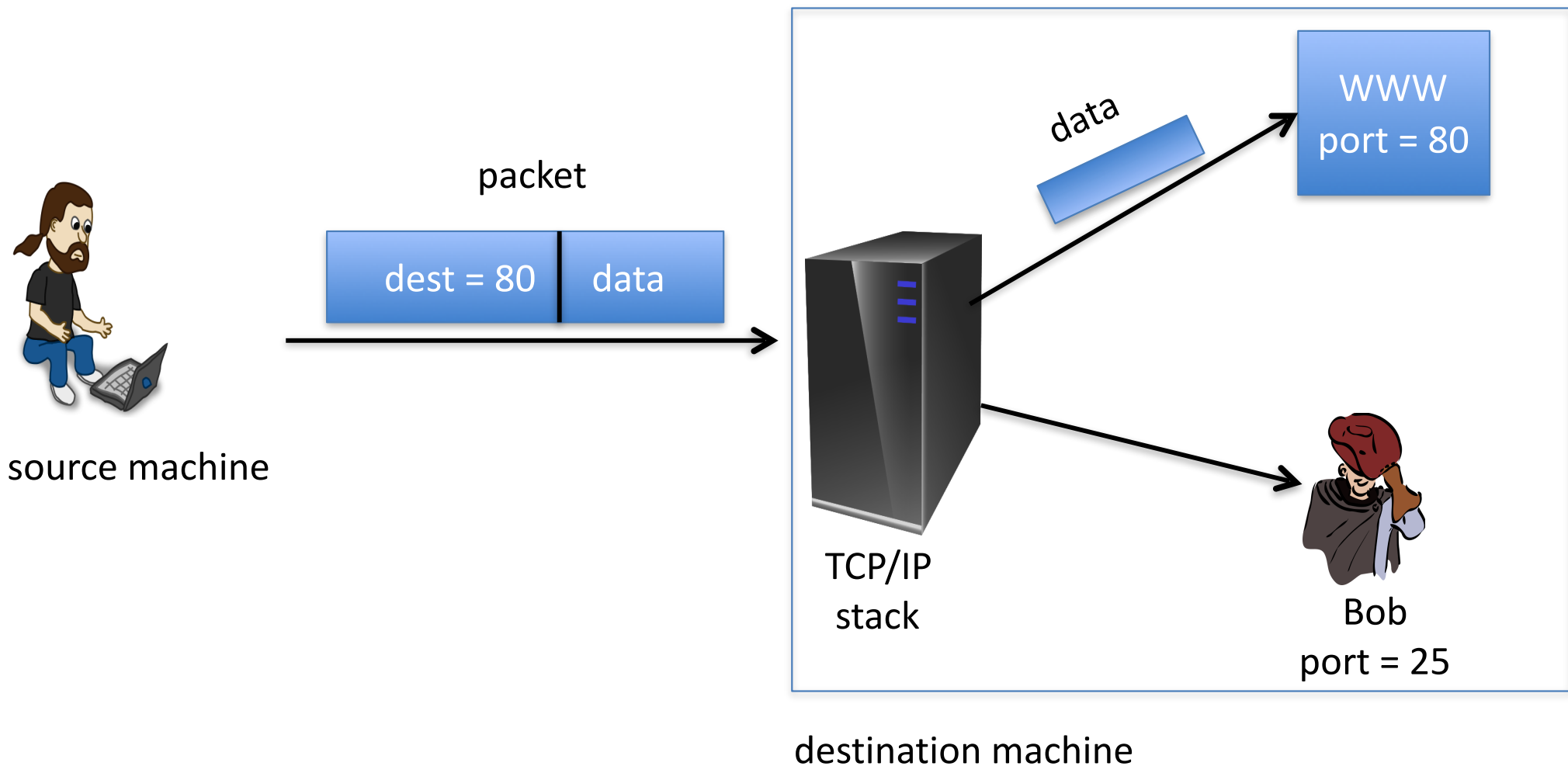
# Story So Far ...

- **Confidentiality:** Secure Encryption (CPA-secure)
  - Single Block messages
  - Multi-block messages
- **Integrity:** Message Authentication Code
  - Using secure block cipher
  - Using collision-resistant hashing

Can we achieve **Confidentiality**  
and **Integrity** at the same time?

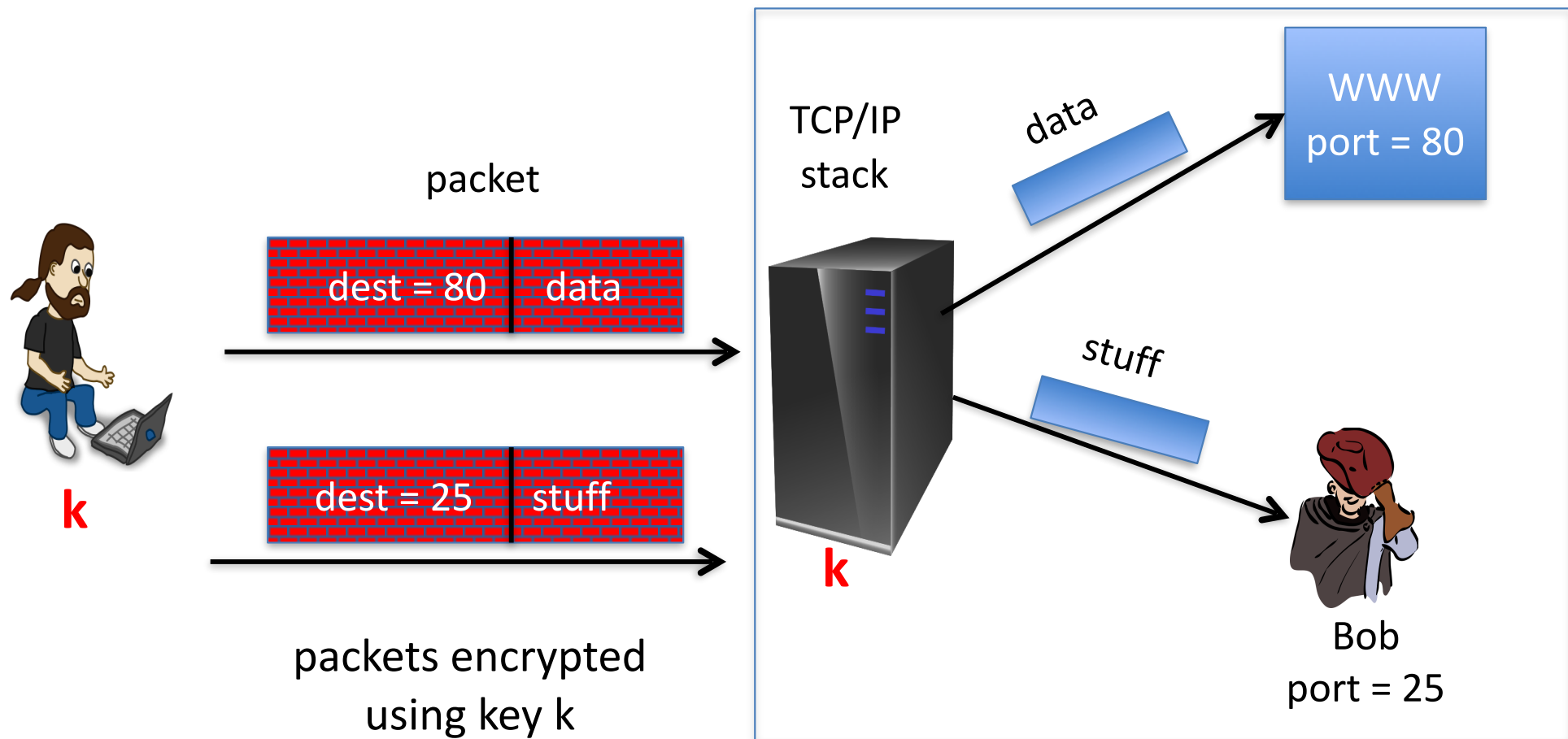
# Sample tampering attacks

TCP/IP: (highly abstracted)



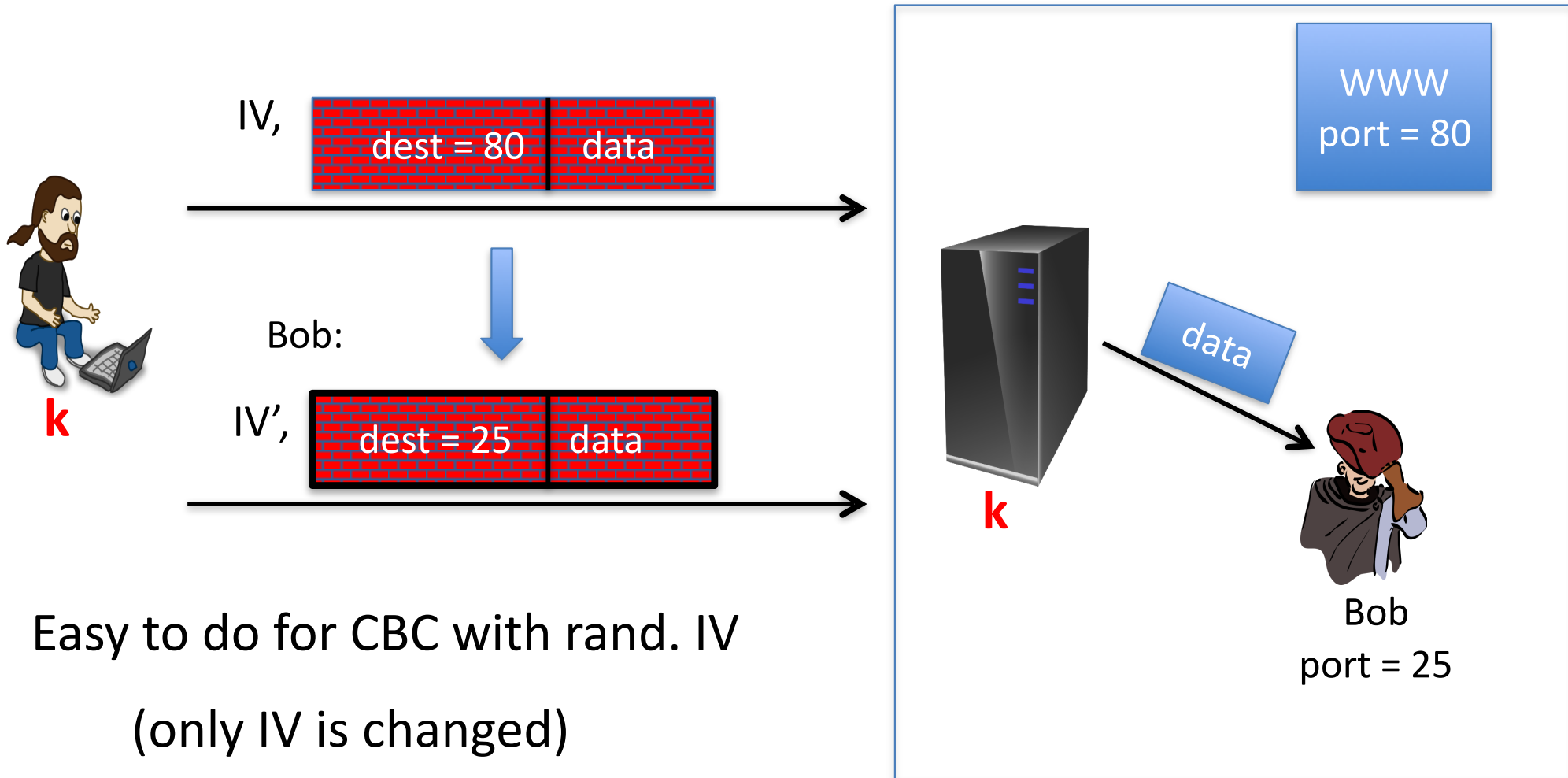
# Sample tampering attacks

IPsec: (highly abstracted)



# Reading someone else's data

Note: attacker obtains decryption of any ciphertext beginning with "dest=25"



Easy to do for CBC with rand. IV  
(only IV is changed)

# The lesson

CPA security cannot guarantee secrecy under active attacks.

Only use one of two modes:


- If message needs integrity but no confidentiality:  
use a **MAC**
- If message needs both integrity and confidentiality:  
use **authenticated encryption** modes

# The Idea

An **authenticated encryption** system (E,D)

As usual:  $E: \mathcal{K} \times \mathcal{M} \rightarrow \mathcal{C}$

but  $D: \mathcal{K} \times \mathcal{C} \rightarrow \mathcal{M} \cup \{\perp\}$

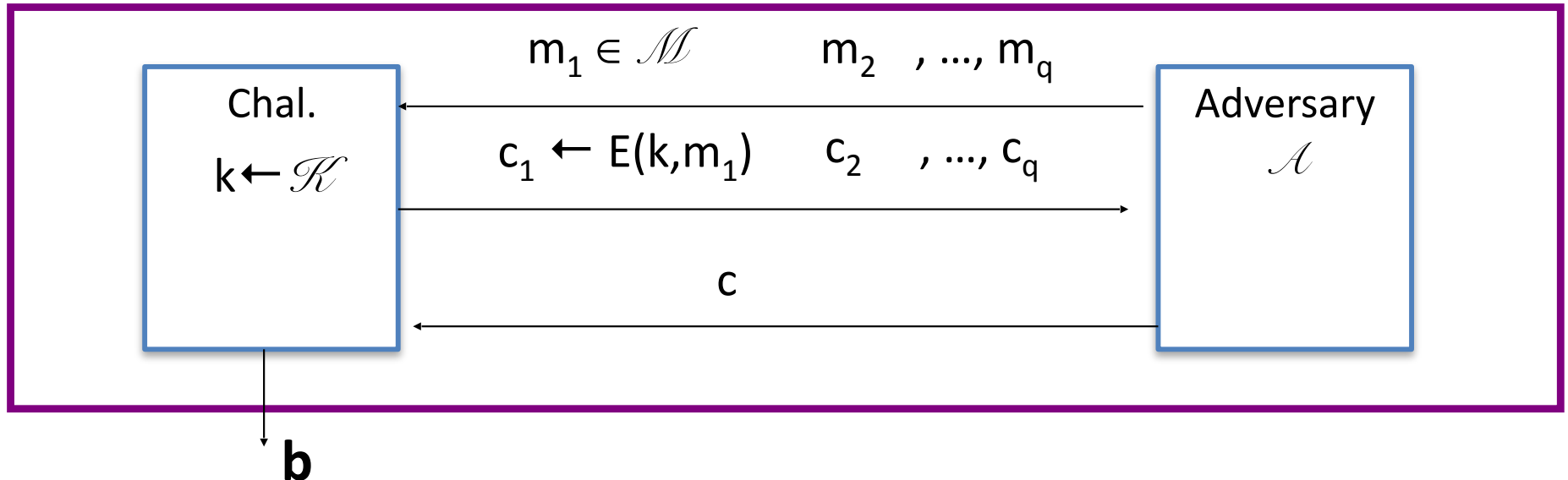
 ciphertext  
is rejected

Security: the system must provide

- semantic security, and
- **ciphertext integrity**:  
attacker cannot create new ciphertexts that  
decrypt properly

# Ciphertext integrity

Let  $(E,D)$  be a cipher with message space  $\mathcal{M}$ .



$$\begin{cases} \mathbf{b}=1 & \text{if } D(k,c) \neq \perp \text{ and } c \notin \{c_1, \dots, c_q\} \\ \mathbf{b}=0 & \text{otherwise} \end{cases}$$

Def:  $(E,D)$  has **ciphertext integrity** if for all “efficient”  $\mathcal{A}$ :

$$\text{Adv}_{\text{CI}}[\mathcal{A}, E] = \Pr[\text{Chal. outputs } 1] \text{ is “negligible.”}$$



# Authenticated Encryption

Def: cipher  $(E,D)$  provides authenticated encryption **(AE)** if it is

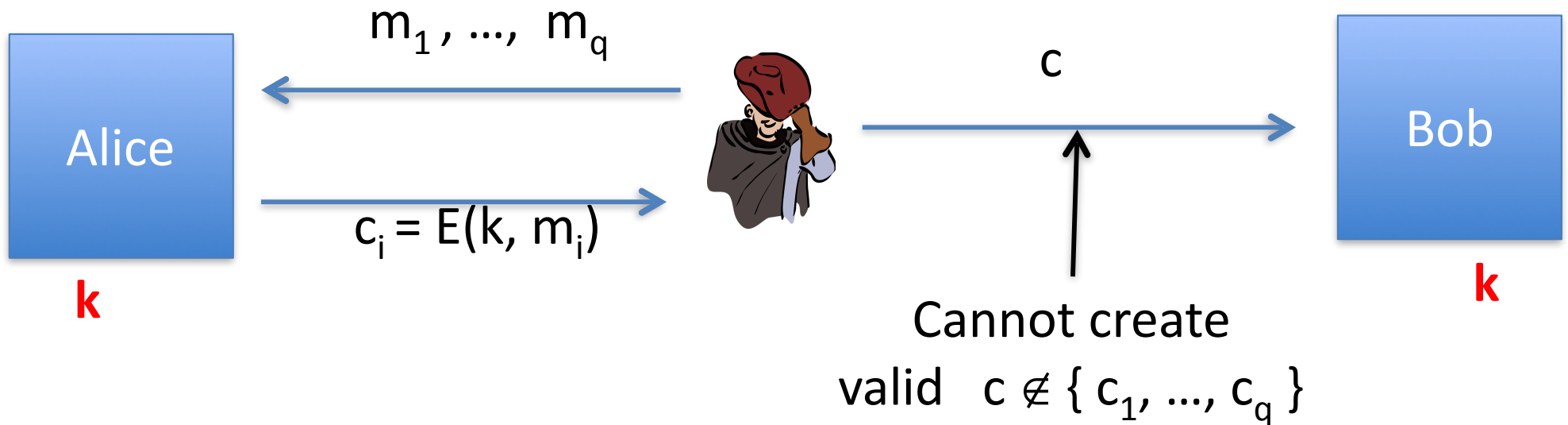
- (1) semantically secure under CPA, and
- (2) has ciphertext integrity

Bad example: CBC with rand. IV does not provide AE

- $D(k, \cdot)$  never outputs  $\perp$ , hence adv. easily wins CI game

# Implication 1: authenticity

Attacker cannot fool Bob into thinking a message was sent from Alice



$\Rightarrow$  if  $D(k, c) \neq \perp$  Bob knows message is from someone who knows  $k$   
(but message could be a replay)

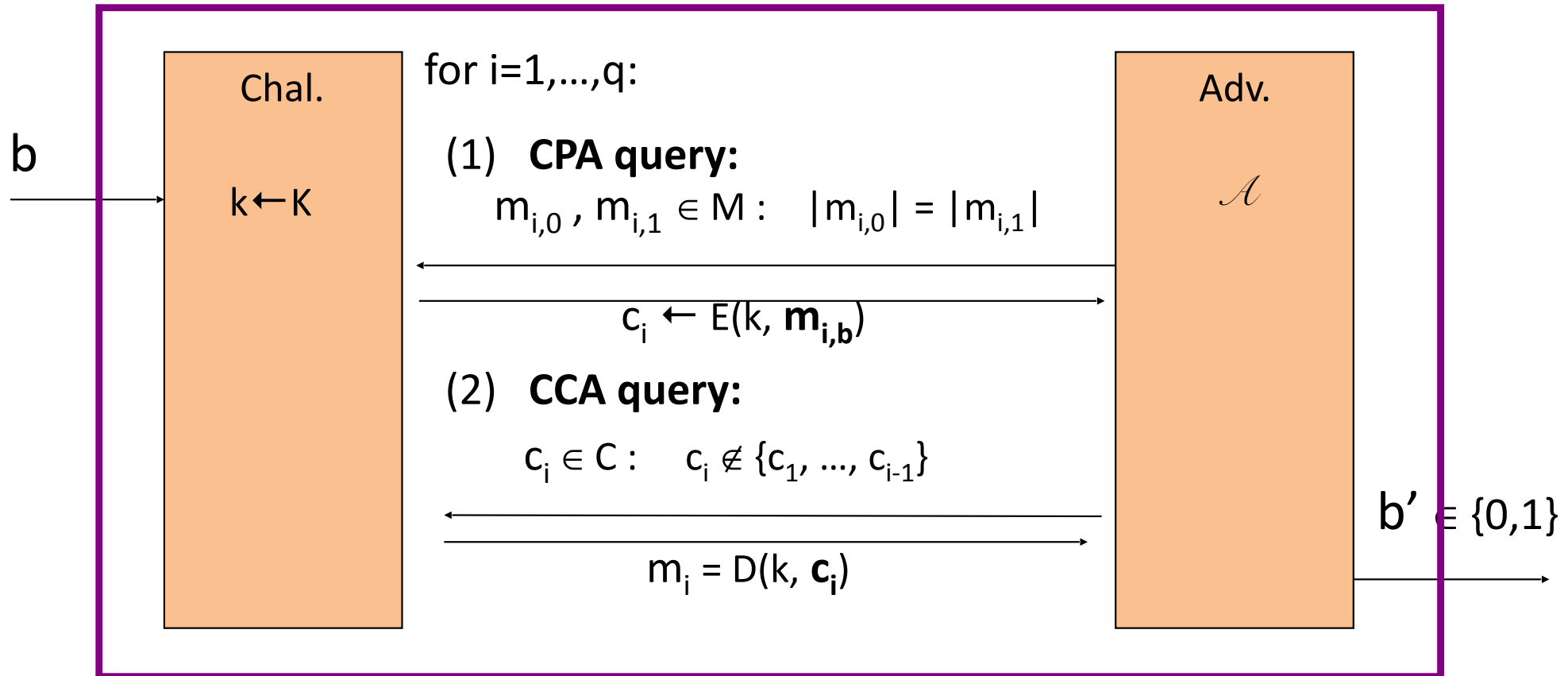
# Implication 2

Authenticated Encryption  $\Rightarrow$

Security against **Chosen Ciphertext Attacks**

# Chosen Ciphertext security -- Definition

For  $b=0,1$  define  $\text{EXP}(b)$ :



$E$  is CCA secure if for all "efficient"  $\mathcal{A}$ :

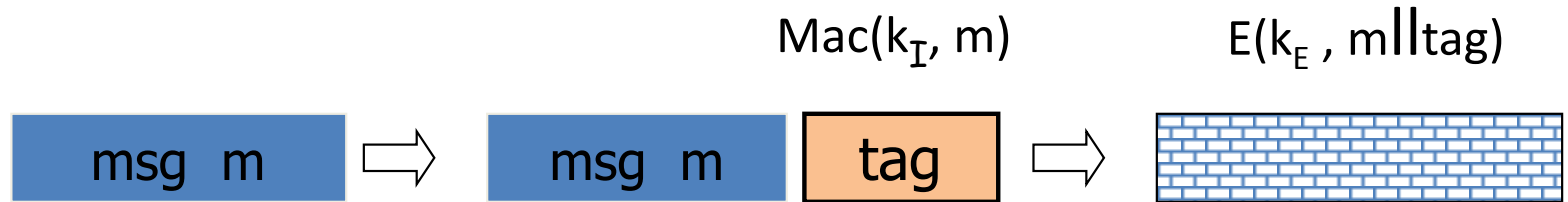
$$\text{Adv}_{\text{CCA}} [A, E] = |\Pr[\text{EXP}(0)=1] - \Pr[\text{EXP}(1)=1]| \text{ is "negligible."}$$

# Combining MAC and Enc (CCA)

Encryption key  $k_E$ .

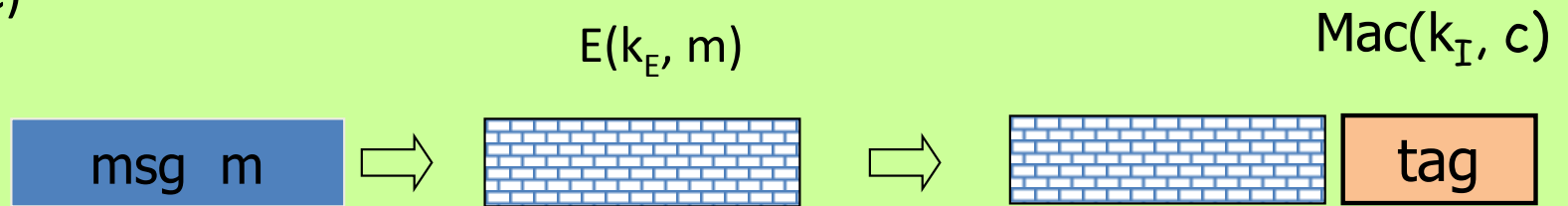
MAC key =  $k_I$

Option 1: (SSL)

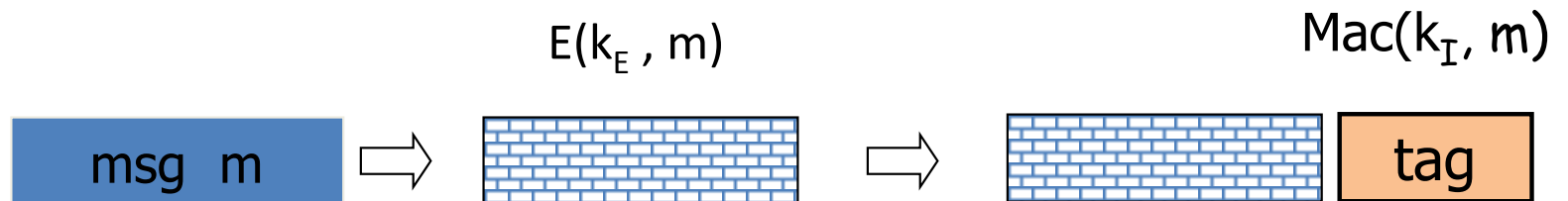


Option 2: (IPsec)

**always  
correct**



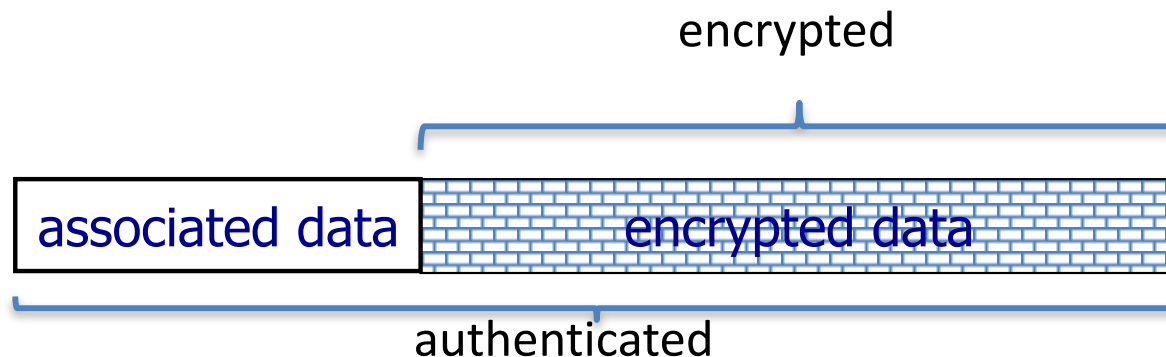
Option 3: (SSH)



# Standards (at a high level)

- **GCM:** CTR mode encryption then CW-MAC  
(accelerated via Intel's PCLMULQDQ instruction)
- **CCM:** CBC-MAC then CTR mode encryption  
(802.11i)
- **EAX:** CTR mode encryption then CMAC

All support AEAD: (auth. enc. with associated data). All are nonce-based.



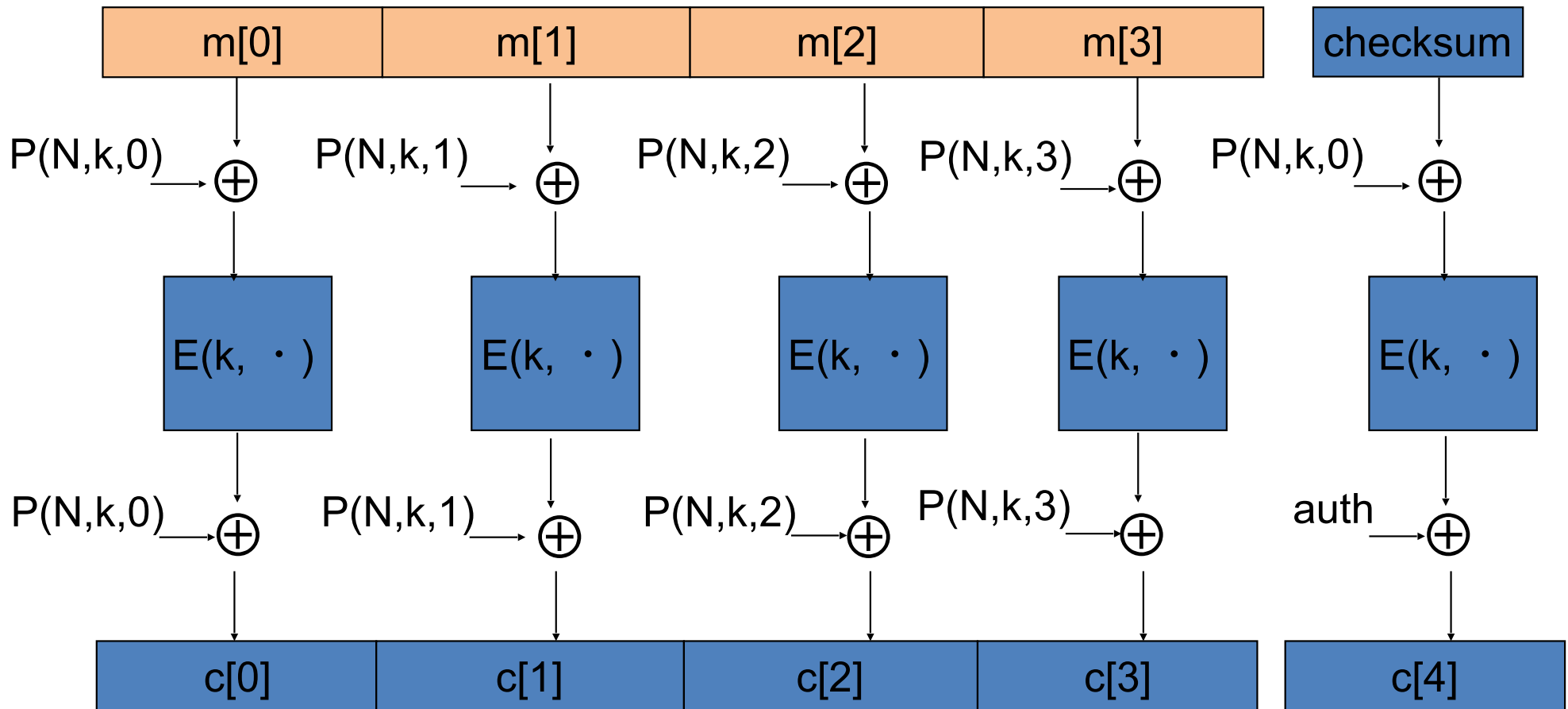
# An example API (OpenSSL)

```
int AES_GCM_Init(AES_GCM_CTX *ain,  
    unsigned char *nonce, unsigned long noncelen,  
    unsigned char *key, unsigned int klen )
```

```
int AES_GCM_EncryptUpdate(AES_GCM_CTX *a,  
    unsigned char *aad, unsigned long aadlen,  
    unsigned char *data, unsigned long datalen,  
    unsigned char *out, unsigned long *outlen)
```

# OCB: a direct construction from a PRP

More efficient authenticated encryption: one  $E()$  op. per block.



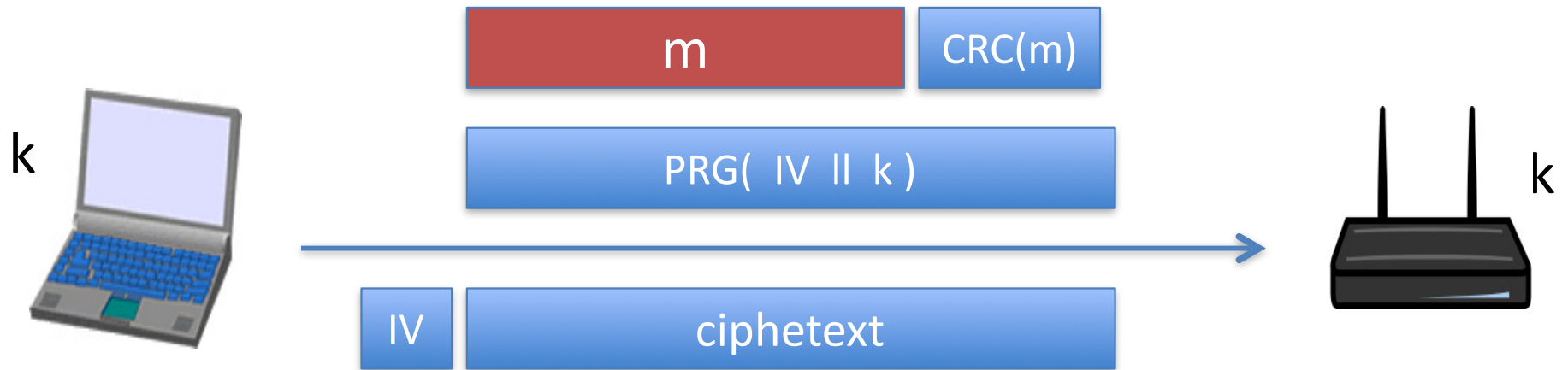
OCB: A Block-Cipher Mode of Operation for Efficient Authenticated Encryption, [RBBK01, CCS]

IETF RFC7253



# 802.11b WEP: how not to do it

## 802.11b WEP:



Previously discussed problems:

two time pad and related PRG seeds

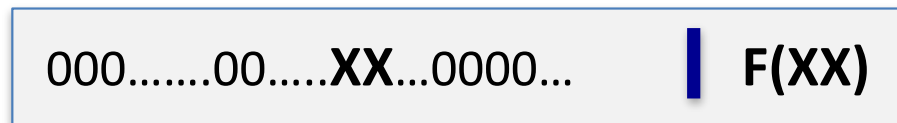
# Active attacks

**Fact:** CRC is linear, i.e.  $\forall m, p: \text{CRC}(m \oplus p) = \text{CRC}(m) \oplus F(p)$

WEP ciphertext:



attacker:



$XX = 25 \oplus 80$



Upon decryption: CRC is valid, but ciphertext is changed !!