Self-applicable probabilistic inference without interpretive overhead

Oleg Kiselyov FNMOC oleg@pobox.com Chung-chieh Shan Rutgers University ccshan@rutgers.edu

16 April 2009

# Patrick Hughes



# Patrick Hughes



#### Probabilistic inference

# $\left. \begin{array}{c} \Pr(W) \\ \Pr(F|W) \\ \text{Observed evidence } F \end{array} \right\} \text{Compute } \Pr(W|F) \text{, etc.}$

	Model (what)	Inference (how)
Toolkit (BNT)	invoke $\rightarrow$	distributions, conditionalization,
Language (BLOG)	random choice, evidence observation,	← interpret

Model (what)

Inference (how)

Toolkituse existing facilities:<br/>libraries, compilers,<br/>types, debuggingadd custom procedures:<br/>just sidestep or extendLanguagesuccinct and natural:<br/>sampling procedures,<br/>relational programscompile models to more<br/>efficient inference code

	Model (what)	Inference (how)
Toolkit (BNT)	use existing facilities: libraries, compilers, types, debugging	add custom procedures: just sidestep or extend
Language (BLOG)	succinct and natural: sampling procedures, relational programs	compile models to more efficient inference code
Today: best of both worlds	invoke $\rightarrow$	← interpret
	models <i>of inference</i> : theory of mind	deterministic parts of models run at full speed

Express both models and inference as programs in the same general-purpose language.

# Outline

#### Expressivity (colored balls) Memoization

#### Inference (music)

Reifying a model into a search tree Importance sampling with look-ahead

#### Self-interpretation (implicature)

Variable elimination Particle filtering Theory of mind

An urn contains an unknown number of balls—say, a number chosen from a [uniform] distribution. Balls are equally likely to be blue or green. We draw some balls from the urn, observing the color of each and replacing it. We cannot tell two identically colored balls apart; furthermore, observed colors are wrong with probability 0.2. How many balls are in the urn? Was the same ball drawn twice? (Milch et al. 2007)

type color = Blue | Green

dist [(0.5, Blue); (0.5, Green)]

type color = Blue | Green

let ball\_color = memo (function b ->
 dist [(0.5, Blue); (0.5, Green)])

type color = Blue | Green

type color = Blue | Green

```
let nballs = 1 + uniform 8 in
let ball_color = memo (function b ->
        dist [(0.5, Blue); (0.5, Green)]) in
let observe = function o ->
    if o <> observed_color(ball_color(uniform nballs))
    then fail ()
```

type color = Blue | Green

```
let nballs = 1 + uniform 8 in
let ball_color = memo (function b ->
        dist [(0.5, Blue); (0.5, Green)]) in
let observe = function o ->
    if o <> observed_color(ball_color(uniform nballs))
    then fail ()
```

```
type color = Blue | Green
let opposite_color = function Blue -> Green
                             | Green -> Blue
let observed_color = function c ->
    dist [(0.8, c); (0.2, opposite_color c)]
 let nballs = 1 + uniform 8 in
 let ball color = memo (function b ->
       dist [(0.5, Blue); (0.5, Green)]) in
 let observe = function o \rightarrow
     if o <> observed_color (ball_color (uniform nballs))
     then fail ()
```

```
let ball_color = memo (function b ->
        dist [(0.5, Blue); (0.5, Green)]) in
let observe = function o ->
    if o <> observed_color(ball_color(uniform nballs))
    then fail ()
```

```
type color = Blue | Green
let opposite_color = function Blue -> Green
                              | Green -> Blue
let observed color = function c \rightarrow
    dist [(0.8, c); (0.2, opposite_color c)]
let model_nballs = function obs () ->
 let nballs = 1 + uniform 8 in
 let ball color = memo (function b ->
       dist [(0.5, Blue); (0.5, Green)]) in
 let observe = function o \rightarrow
     if o <> observed_color(ball_color(uniform nballs))
     then fail () in
 Array.iter observe obs; nballs
```

```
type color = Blue | Green
let opposite_color = function Blue -> Green
                             | Green -> Blue
let observed_color = function c ->
    dist [(0.8, c); (0.2, opposite_color c)]
let model_nballs = function obs () ->
 let nballs = 1 + uniform 8 in
 let ball color = memo (function b ->
       dist [(0.5, Blue); (0.5, Green)]) in
 let observe = function o \rightarrow
     if o <> observed_color(ball_color(uniform nballs))
     then fail () in
 Array.iter observe obs; nballs
normalize (sample_reify 17 10000 (model_nballs
```

```
type color = Blue | Green
let opposite_color = function Blue -> Green
                             | Green -> Blue
let observed_color = function c ->
    dist [(0.8, c); (0.2, opposite_color c)]
let model_nballs = function obs () ->
 let nballs = 1 + uniform 8 in
let ball_color = memo (function b ->
       dist [(0.5, Blue); (0.5, Green)]) in
 let observe = function o \rightarrow
     if o <> observed_color(ball_color(uniform nballs))
     then fail () in
 Array.iter observe obs; nballs
normalize (sample_reify 17 10000 (model_nballs
```

#### Outline

#### Expressivity (colored balls) Memoization

#### ► Inference (music)

Reifying a model into a search tree Importance sampling with look-ahead

Self-interpretation (implicature)

Variable elimination Particle filtering Theory of mind

С

```
type 'a vc = V of 'a | C of (unit -> 'a pV)
and 'a pV = (float * 'a vc) list
```



```
type 'a vc = V of 'a | C of (unit -> 'a pV)
and 'a pV = (float * 'a vc) list
```



type 'a vc = V of 'a | C of (unit -> 'a pV)
and 'a pV = (float \* 'a vc) list



type 'a vc = V of 'a | C of (unit -> 'a pV)
and 'a pV = (float \* 'a vc) list

Depth-first traversal is exact inference by brute-force enumeration.



type 'a vc = V of 'a | C of (unit -> 'a pV)
and 'a pV = (float \* 'a vc) list

Inference procedures cannot access models' source code.



- Implementedby representing(Filinski 1994)a state monad transformer(Moggi 1990)applied to a probability monad(Giry 1982)using shift and reset(Danvy & Filinski 1989)to operate on first-class(Felleisen et al. 1987)delimited continuations(Strachey & Wadsworth 1974)
  - model runs inside reset (like an exception handler)
  - dist and fail perform shift (like throwing an exception)
  - memo mutates thread-local storage

C Probability mass  $p_c = 1$ 



Probability mass  $p_c = 1$ 

1. Expand one level.



Probability mass  $p_c = 1$  (.2, Green)

- 1. Expand one level.
- 2. Report shallow successes.



Probability mass  $p_c = .75$  (.2, Green)

- 1. Expand one level.
- 2. Report shallow successes.
- 3. Expand one more level and tally open probability.



Probability mass  $p_c = .75$  (.2, Green)

- 1. Expand one level.
- 2. Report shallow successes.
- 3. Expand one more level and tally open probability.
- 4. Randomly choose a branch and go back to 2.



Probability mass  $p_c = .75$ (.2, Green) (.6, Blue)

- 1. Expand one level.
- 2. Report shallow successes.
- 3. Expand one more level and tally open probability.
- 4. Randomly choose a branch and go back to 2.



Probability mass  $p_c = 0$ (.2, Green) (.6, Blue)

- 1. Expand one level.
- 2. Report shallow successes.
- 3. Expand one more level and tally open probability.
- 4. Randomly choose a branch and go back to 2.



Probability mass  $p_c = 0$  (.2, Green) (.6, Blue)

- 1. Expand one level.
- 2. Report shallow successes.
- 3. Expand one more level and tally open probability.
- 4. Randomly choose a branch and go back to 2.

#### Music model

Pfeffer's test of importance sampling (2007): motivic development in early Beethoven piano sonatas



Want  $Pr(D = \cdots | S = \cdots)$ . Exact inference and rejection sampling are infeasible. Implemented using lists with stochastic parts.

# Typical inference results



#### Outline

#### Expressivity (colored balls) Memoization

#### Inference (music)

Reifying a model into a search tree Importance sampling with look-ahead

#### Self-interpretation (implicature)

Variable elimination Particle filtering Theory of mind

#### Models of inference

Inference procedures and models

- are written in the same general-purpose language
- use the same stochastic primitive dist

#### Models of inference

Inference procedures and models

- are written in the same general-purpose language
- use the same stochastic primitive dist

so inference procedures can be invoked by models

inference (function ()  $\rightarrow$ 

... inference (function ()  $\rightarrow$  ...) ...)

and deterministic parts run at full speed.

Program generation with mutable state and control effects.

#### Models of inference

Inference procedures and models

- are written in the same general-purpose language
- use the same stochastic primitive dist

so inference procedures can be invoked by models

inference (function ()  $\rightarrow$ 

... inference (function ()  $\rightarrow$  ...) ...)

and deterministic parts run at full speed.

Program generation with mutable state and control effects.

One common usage pattern: reify-infer-reflect

- Brute-force enumeration becomes variable elimination
- Sampling becomes particle filtering

# Theory of mind

Instances abound:

- False-belief (Sally-Anne) task
- Trading securities
- Teacher's hint to student
- Gricean reasoning in language use

# Theory of mind

Instances abound:

- False-belief (Sally-Anne) task
- Trading securities
- Teacher's hint to student
- Gricean reasoning in language use
  - 1. "Some professors are coming to the party."
  - 2. "All professors are coming to the party."
  - 3. "Some but not all professors are coming to the party."

Trade-off between precision and ease of comprehension?

# Theory of mind

Instances abound:

- False-belief (Sally-Anne) task
- Trading securities
- Teacher's hint to student
- Gricean reasoning in language use
  - 1. "Some professors are coming to the party."
  - 2. "All professors are coming to the party."
  - 3. "Some but not all professors are coming to the party."

Trade-off between precision and ease of comprehension?

Crucial for collaboration among human and computer agents!

Want executable models.

A bounded-rational agent's theory of bounded-rational mind  $\sim$  approximate inference about approximate inference

#### Marr's computational vs algorithmic models



#### world $W \in \{0 \text{ come, } 1 \text{ come, } 2 \text{ come, } 3 \text{ come}\} \times \cdots$

- action  $A \in \{ \text{feed 0}, \text{feed 1}, \text{feed 2}, \text{feed 3} \}$
- form  $F \subseteq \{\text{some, all, no, not all}\}$

#### Marr's computational vs algorithmic models



world  $W \in \{0 \text{ come, } 1 \text{ come, } 2 \text{ come, } 3 \text{ come}\} \times \cdots$ 

- action  $A \in \{ \text{feed 0}, \text{feed 1}, \text{feed 2}, \text{feed 3} \}$
- form  $F \subseteq \{\text{some, all, no, not all}\}$

#### Marr's computational vs algorithmic models



form  $F \subseteq \{\text{some, all, no, not all}\}$ 

#### Marr's comm

#### inference $\Pr(F)$

**model** Pr(W), U(A, W)

**inference** Pr(A|F, true)

**model** Pr(W), Pr(true|W, F), U(A, W)





A computational model of the modeler nests an algorithmic model of the modelee: invoke inference recursively, without interpretive overhead.

#### Summary

# Express both models and inference as programs in the same general-purpose language.

- Combine strengths of toolkits and standalone languages
- Deterministic parts of models run at full speed
- Models can invoke inference without interpretive overhead
- Theory of mind: inference about approximate inference
- A variety of inference methods: variable elimination, particle filtering, importance sampling, ...?