# Quotation and effects in natural language Three applications

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Computational Linguistics





?

#### **Outline**

#### ► Past: Mixed quotation

Present: Quantifier scope

Quotation for programming: code generation Control for programming: let insertion Control for linguistics: quantification Quotation for linguistics: inverse scope

Future: Rational metaprogramming

# Anaphora as state

Apparently, the idea of **meeting participants** making **their** own reservations at the hotel does not work well for them.

# Mixed quotation

"Bachelor" has eight letters.

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#### Direct

Quine said, "Quotation has a certain anomalous feature".

#### Indirect

Quine said that quotation has a certain anomalous feature.

#### Mixed

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# Mixed quotation

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#### Direct

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#### Indirect

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#### Mixed

Quine said that quotation "has a certain anomalous feature".

Bush said he has an "ecelectic" reading list.

Bush said the enemy "misunderestimates me".

# Anaphora in quotation

The professor said she requires "every student in my class who lives on campus" to bring their homework to her office.

#### Professor to journalist:

I require every student in my class who lives on campus to drop their work into this box.

Run with state?

# Anaphora in quotation

The professor said she requires "every student in my class who lives on campus" to bring their homework to her office.

#### Professor to journalist:

I require every student in my class who lives on campus to drop their work into this box.

Run with state?

\* The professor told every student in her class who lives on campus to "bring their homework to my office".

#### Professor to John:

Please bring your Lordship's homework to my office.

#### Professor to Mary:

Please bring your Ladyship's homework to my office.

No cross-stage persistence?

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# Staging power

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```
let rec power1 x = function
     I 0 -> 1
     \mid n \rightarrow x * power1 x (pred n)
▶ val power1: int -> int -> int = <fun>
let test1 = power1 2 11
▶ val test1: int = 2048
let rec power2 x = function
     \mid 0 \rightarrow \langle 1 \rangle
     \mid n \rightarrow \langle x * (power2 x (pred n)) \rangle
▶ val power2: ('a,int) code -> int -> ('a,int) code = <fun>
let test2 = \langle \text{fun x} \rightarrow \text{(power2 } \langle \text{x} \rangle 11) \rangle
```

# Interpreting English

```
John loves Mary
-: bool = true

John loves himself
-: bool = false

Someone loves John
-: bool = true
```

# Interpreting English

```
John loves Mary
-: bool = true
John loves himself
▶ -: bool = false
Someone loves John
▶ -: bool = true
John loves Mary
▶ -: ('a,bool) code = ⟨love Mary John⟩
John loves himself
▶ -: ('a,bool) code = ⟨love John John⟩
Someone loves John
▶ -: ('a,bool) code = ⟨List.exists people (love John)⟩
```

#### The need to insert let

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# Inserting let in continuation-passing (or monadic) style

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```
let square5 x k = \langle let r = x * x in (k \langle r \rangle) \rangle
```

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```
let square5 x k = \langle let r = x * x in (k \langle r \rangle) \rangle
let rec power5 x k = function
      | 0 -> k \langle 1 \rangle
      | 1 -> k x
      1 \text{ n when n mod } 2 = 0
            \rightarrow square5 x (fun s \rightarrow power5 s k (n/2))
      \mid n -> square5 x (fun s -> power5 s (fun r ->
                                                   k \langle r * x \rangle
                                                (n/2)
let test5 = \langle \text{fun x} \rightarrow \text{``(power5 } \langle \text{x} \rangle \text{ (fun r} \rightarrow \text{r)} 11) \rangle
▶ val test5: ('a, int -> int) code
   = \langle \text{fun } x \rightarrow \text{let } r1 = x * x \text{ in} \rangle
                     let r2 = r1 * r1 in
                     let r3 = r2 * r2 in (r3 * r1) * x
```

# Inserting let in direct style

```
let square6 x = shift (fun k -> \langle let r = x * x in (k \langle r \rangle) \rangle)
let rec power6 x = function
      | 0 -> \(1\)
      | 1 -> x
      | n when n mod 2 = 0 \rightarrow power6 (square6 x) (n/2)
      | n \rightarrow \langle (power6 (square6 x) (n/2)) * x \rangle
let test6 = \langle \text{fun } x \rightarrow \text{``(reset (fun () -> power6 } \langle x \rangle 11))} \rangle
▶ val test6: ('a, int -> int) code
   = \langle \text{fun } x \rightarrow \text{let } r1 = x * x \text{ in} \rangle
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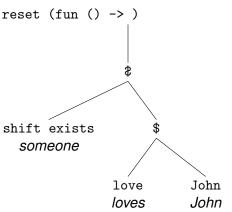
# Shifting gears

```
type person = John | Mary
let people = [John; Mary]
let love (x: person) (y: person) = x != y
let x ? f = f x
      John & (love $ Mary)
         John loves Mary.
                              John
                              John
                                 love
                                        Mary
                                 loves
                                        Mary
```

# In-situ quantification

```
let forall (f: person -> bool) = List.for_all f people
let exists (f: person -> bool) = List.exists f people
```

#### Someone loves John.



### In-situ quantification

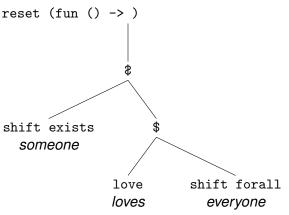
```
let forall (f: person -> bool) = List.for_all f people
let exists (f: person -> bool) = List.exists f people
```

# Mary loves everyone. reset (fun () -> ) Mary Mary shift forall love loves everyone

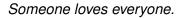
# In-situ quantification

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let forall (f: person -> bool) = List.for_all f people
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### Someone loves everyone.

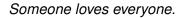


Someone loves everyone.



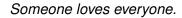


Children...
There's a time and a place for everything





Children...
There's a time and a place for everything, and it's called college.





Children...

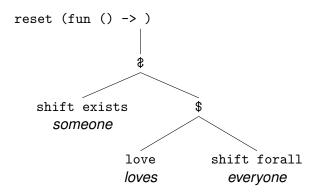
There's a time and a place for everything, and it's called college.

Require left-to-right evaluation for other side effects:

- \* His mother loves everyone.
- \* What did who buy?
- \* Anyone loves no one.

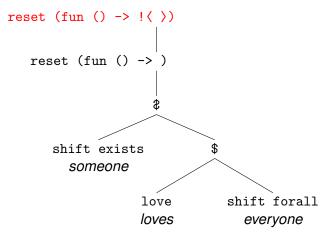
# Inverse scope as quotation

Someone loves everyone.



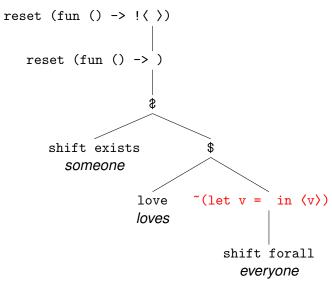
# Inverse scope as quotation

"Someone loves everyone".



# Inverse scope as quotation

"Someone loves" everyone.



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► Future: Rational metaprogramming

# Speaker and hearer model each other

Isn't it getting chilly in here?

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Isn't it getting chilly in here?

A hotel cleaner enters a room and starts to clean it. A female guest emerges from the shower. The cleaner says "Excuse me sir" and exits.

# Rational metaprogramming

To model the beliefs, desires, and intentions of agents who have beliefs about each other's intentions, about each other's desires about each other's beliefs, and so on,

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To model the beliefs, desires, and intentions of agents who have
    beliefs about each other's intentions,
            about each other's desires about each other's beliefs.
                                                       and so on,
we model
  intentions to perform actions as programs.
     beliefs as probability distributions.
             (weighted nondeterminism → stochastic programs)
    desires as utility functions.
             (rational choice → rational programs)
```

# Rational metaprogramming

To model the beliefs, desires, and intentions of agents who have beliefs about each other's intentions,

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#### we model

```
intentions to perform actions as programs.
```

beliefs as probability distributions.

(weighted nondeterminism → stochastic programs)

desires as utility functions.

(rational choice → rational programs)

One agent's model of another is a probability distribution over (quoted) rational programs.

We need a modal type system and efficient self-interpretation.

# Summary

Quotation goes well with effects (state, control, nondeterminism), so that code does not have to be generated in lexical order.

But we want a type system that prevents scope extrusion.

Multigrained theories of quotation: the less intensional a theory, the more cross-stage persistence it allows?

Levels of quotation are not quite levels of control operators.

#### Reckless let insertion

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```
let test6 = \langle \text{fun x} - \rangle ~(reset (fun () ->
      power6 \langle x \rangle 11))
▶ val test6: ('a, int -> int) code
   = \langle \text{fun } x \rightarrow \text{let } r1 = x * x \text{ in} \rangle
                     let r2 = r1 * r1 in
                     let r3 = r2 * r2 in (r3 * r1) * x
let test7a = \langle \text{fun x} - \rangle ~(reset (fun () ->
      \langle \text{let } y = x + 1 \text{ in } (\text{power6 } \langle y \rangle 11) \rangle \rangle \rangle
▶ val test7a: ('a, int -> int) code
   = \langle \text{fun } x \rightarrow \text{let } r1 = y * y \text{ in} \rangle
                     let r2 = r1 * r1 in
                     let r3 = r2 * r2 in
                     let y = x + 1 in (r3 * r1) * y
```

#### Reckless let insertion

```
let test6 = \langle \text{fun x} - \rangle ~(reset (fun () ->
      power6 \langle x \rangle 11))\rangle
▶ val test6: ('a, int -> int) code
   = \langle \text{fun } x \rightarrow \text{let } r1 = x * x \text{ in} \rangle
                      let r2 = r1 * r1 in
                      let r3 = r2 * r2 in (r3 * r1) * x
let test7b = \langle \text{fun x} - \rangle ~(reset (fun () ->
      \langle \text{let } y = x + 1 \text{ in } (\text{reset } (\text{fun } () \rightarrow \text{power6 } \langle y \rangle 11)) \rangle \rangle
▶ val test7b: ('a, int -> int) code
   = \langle \text{fun } x \rightarrow \text{let } y = x + 1 \text{ in} \rangle
                      let r1 = y * y in
                      let r2 = r1 * r1 in
                      let r3 = r2 * r2 in (r3 * r1) * y
```